Molecules as Automata Representing Biochemical Systems as Collectives of Interacting Automata

Luca Cardelli

Microsoft Research

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http://LucaCardelli.name

Engineering Method



Direct Engineering (Synthetic Biology)

Scientific Method



Reverse Engineering (Systems Biology)

Engineering Method



Direct Engineering

Scientific Method



Reverse Engineering









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Motivation: Cells Compute

- No survival without computation!
 - Finding food
 - Avoiding predators
- How do they compute?
 - Unusual computational paradigms.
 - Proteins: do they work like electronic circuits?
 - Genes: what kind of software is that?
- Signaling networks
 - Clearly "information processing"
 - They are "just chemistry": molecule interactions
 - But what are their principles and algorithms?
- Complex, higher-order interactions
 - MAPKKK = MAP Kinase Kinase Kinase: that which operates on that which operates on that which operates on protein.
- General models of biological computation
 - What are the appropriate ones?



Discrete State Transitions

Epidermal Growth Factor Receptor Pathway Map

Kanase Cida (177), Turciso Manazoka (4, Pendadi Ratano (1774) (5 Te tysee long milia, (5 Center a futures long and langing contemp),



Modeling Approach

- We believe that {petri nets, process algebra, term rewriting, multiagent systems} are {better, complementary} for modeling biological systems than {SBML, Kohn charts, chemical reactions, ODEs}.
- We take a paper from the literature (usually ODEs or chemical reactions) and "code it up" in e.g. Petri nets.
- How do we know that's the "same system"? How do we convince mathematical biologists that we are doing the "right thing"?

(Macro-) Molecules as (Interacting) Automata

Process Algebra

- Reactive systems (living organisms, computer networks, operating systems, ...)
 - Math is based on *entities that react/interact with their environment* ("*processes"*), not on functions from domains to codomains.
- Concurrent
 - Events (reactions/interactions) happen concurrently and asynchronously, not sequentially like in function composition.
- Stochastic
 - Or probabilistic, or nondeterministic, but is never about deterministic system evolution.
- Stateful
 - Each concurrent activity ("process") maintains its own local state, as opposed to stateless functions from inputs to outputs.
- Discrete
 - Evolution through discrete transitions between discrete states, not incremental changes of continuous quantities.
- Kinetics of interaction
 - An "interaction" is anything that moves a system from one state to another.



Kinetic laws:

Kinetic laws:

Two complementary actions may result in an interaction.

Kinetic laws:

Two complementary actions may result in an interaction. A decay may happen spontaneously.

Interactions have rates. Actions DO NOT have rates.

The equivalent process algebra model

Interactions in a Population

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Interactions in a Population

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Interactions in a Population

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Interactions in a Population (2)

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Interactions in a Population (2)

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CTMC Semantics

CTMC

Reactions vs. Components

Groupies and Celebrities

Groupies and Celebrities

A stochastic collective of celebrities:

Stable because as soon as a A finds itself in the majority, it is more likely to find somebody in the same state, and hence change, so the majority is weakened.

Groupies and Celebrities

Groupie									
(wants	to	be	like	som	ebody	differ	rent)		

directive sample 1.0 1000	a@1.0
new a@1.0:chan()	b@1.0
new b@1.0:chan()	
let A() = do !a; A() or ?b; B() and B() = do !b; B() or ?a; A()	
run 100 of (A() B())	

A stochastic collective of groupies:

Unstable because within an A majority, an A has difficulty finding a B to emulate, but the few B's have plenty of A's to emulate, so the majority may switch to B. Leads to deadlock when everybody is in the same state and there is nobody different to emulate.

Both Together

A way to break the deadlocks: Groupies with just a few Celebrities

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Hysteric Groupies

We can get more regular behavior from groupies if they "need more convincing", or "hysteresis" (history-dependence), to switch states.

Some Devices

Ultrasensitive Switch

Cascade Amplifier

Symmetric Wave Generator

More Devices

A0

В0 С0

Oscillator

Repressilator (1 of 3 similar gates)

Semantics of Collective Behavior

The Two Semantic Sides of Chemistry

These diagrams commute via appropriate maps.

L. Cardelli: "On Process Rate Semantics" (TCS)

L. Cardelli: "A Process Algebra Master Equation" (QEST'07)

Quantitative Process Semantics

From CGF to Chemistry

Chemical Reactions

Unary Reactiond[A]/dt = -r[A]Exponential DecayHetero Reaction $d[A_i]/dt = -r[A_1][A_2]$ Mass Action LawHomeo Reaction $d[A]/dt = -2r[A]^2$ Mass Action Law(assuming $A \neq B_i \neq A_j$ for all i,j)

No other reactions!

URNAL OF CHEMICAL PHYSICS VOLUME 113, NUMBER 1 he chemical Langevin equation Daniel T. Gillespie ³⁾ Research Department, Code 4T4100D, Naval Air Warfare Center, China Lake, California 93555 Genuinely trimolecular reactions do not physically occur in dilute fluids with any appreciable frequency. Apparently timolecular reactions in a fluid are usually the combined esult of two bimolecular reactions and one monomolecular eaction, and involve an additional short-lived species.		Chapter IV: Ch [David A. Reckho reactions may elementary. <u>Elem</u> reactions that oc written, without These reactions of or two reactants, involve a series o reactions. Many of reactions are nor reactions with an greater than two integer reaction	emical Kinetics w, CEE 572 Course] be either elementary or non- entary reactions are those cour exactly as they are any intermediate steps. almost always involve just one Non-elementary reactions f two or more elementary complex environmental n-elementary. In general, overall reaction order , or reactions with some non- order are non-elementary.	THE COLLISION THEORY OF REACTION RATES www.chemguide.co.uk The chances of all this happening if your reaction needed a collision involving more than 2 particles are remote. All three (or more) particles would have to arrive at exactly the same point in space at the same time, with everything lined up exactly right, and having enough energy to react. That's not likely to happen very often!
Trimolecular rea $A + B + C \rightarrow^{r} D$ the measured "r aggregate of e.g $A + B \leftrightarrow AB$ $AB + C \rightarrow D$	ctions:) " is an (imperfect) .:	Er th (a	s $__\r P$ e "r" is given by Mi pproximated stead E + S \leftrightarrow ES ES \rightarrow P + E	chaelis-Menten y-state) laws:

Chemical Ground Form (CGF)



$$A = !a_{(r)}; A \oplus ?a_{(r)}; B$$
$$B = ?a_{(r)}; A \oplus \tau_{(s)}; A$$



$A = !a; A \oplus ?a; B$ $B = ?a; A \oplus \tau_{(s)}; A$



$A = !a; A \oplus ?a; B$ $B = ?a; A \oplus$ $\tau_{(s)}; A$



Interacting Automata	 Discrete Chemistry
initial states A A A	initial quantities #A ₀
A @r A'	A→r A'
A ?a A' B !a' [@] r B'	A+B → r A'+B'
?a A !a A' @r A"	A+A ⊶•²r A'+A″



From CGF to Chemistry: Ch(E)

E ::= 0 : X=M, E	Reagents
M ::= Ο ∶ π;Ρ ⊕ M	Molecules
P ::= O : X P	Solutions
$\pi ::= \tau_{(r)} : ?a_{(r)} : !a_{(r)}$	Interactions (delay, input, output)
<i>CG</i> F ::= E,P	Reagents plus Initial Conditions

E.X.i ≝ the i-th ⊕-summand of the molecule M associated with the X reagent of E

Chemical reactions for E,P:

(N.B.: <...> are reaction tags to obtain multiplicity of reactions, and P is P with all the | changed to +)

 $\begin{array}{l} Ch(E) := \\ \{(<X.i>: X \to^{r} P) \text{ s.t. E.X.i} = \tau_{(r)}; P\} \cup \\ \{(<X.i,Y.j>: X + Y \to^{r} P + Q) \text{ s.t. } X \neq Y, E.X.i = ?a_{(r)}; P, E.Y.j = !a_{(r)}; Q\} \cup \\ \{(<X.i,X.j>: X + X \to^{2r} P + Q) \text{ s.t. E.X.i} = ?a_{(r)}; P, E.X.j = !a_{(r)}; Q) \rangle \in E\} \end{array}$

Initial conditions for P:



From Chemistry to CGF

From Chemistry to CGF (by example)

x: $B \rightarrow^{s} A$ b: $A+B \rightarrow^{r} A+A$ c: $A+A \rightarrow^{2r} A+B$ Unique reaction names



From Chemistry to CGF (by example)

- $\mathbf{x}: \mathbf{B} \to^{\mathsf{s}} \mathbf{A}$
- b: $A+B \rightarrow^{r} A+A$ c: $A+A \rightarrow^{2r} A+B$

	$x_{(s)}$	b _(r)	c _(r)
A			
В	τ; Α		

- 1: Fill the matrix by columns:
 - Degradation reaction $v_i: X \rightarrow k_i P_i$ add $t; P_i$ to $\langle X, v_i \rangle$.

- $\begin{array}{ll} \mathsf{x:} & \mathsf{B} \to^{\mathsf{s}} \mathsf{A} \\ \mathsf{b:} & \mathsf{A} + \mathsf{B} \to^{\mathsf{r}} \mathsf{A} + \mathsf{A} \end{array} \end{array}$
- c: $A+A \rightarrow^{2r} A+B$

	$x_{(s)}$	b _(r)	C _(r)
A		?;A A	
В	τ; Α	!;0	

1: Fill the matrix by columns:

Degradation reaction $v_i: X \rightarrow k_i P_i$ add $t;P_i$ to $\langle X, v_i \rangle$. Hetero reaction $v_i: X+Y \rightarrow k_i P_i$ add $2;P_i$ to $\langle X, v_i \rangle$ and 1;0 to $\langle Y, v_i \rangle$

- $\begin{array}{ll} \mathbf{x} & \mathbf{B} \to^{\mathsf{s}} \mathbf{A} \\ \mathbf{b} & \mathbf{A} + \mathbf{B} \to^{\mathsf{r}} \mathbf{A} + \mathbf{A} \end{array}$
- c: $A+A \rightarrow^{2r} A+B$

	$x_{(s)}$	b _(r)	C _(r)
A		?;A A	?;A B !;0
В	τ; Α	!;0	

1: Fill the matrix by columns:

Degradation reaction $v_i: X \rightarrow k_i P_i$ add $t;P_i$ to $\langle X, v_i \rangle$. Hetero reaction $v_i: X+Y \rightarrow k_i P_i$ add $2;P_i$ to $\langle X, v_i \rangle$ and 1;0 to $\langle Y, v_i \rangle$ Homeo reaction $v_i: X+X \rightarrow k_i P_i$ add $2;P_i$ and 1;0 to $\langle X, v_i \rangle$

- x: B →^s A b: A+B →^r A+A c: A+A →^{2r} A+B
- 1: Fill the matrix by columns:

Degradation reaction $v_i: X \rightarrow k_i P_i$ add $t;P_i$ to $\langle X, v_i \rangle$. Hetero reaction $v_i: X+Y \rightarrow k_i P_i$ add $2;P_i$ to $\langle X, v_i \rangle$ and 1;0 to $\langle Y, v_i \rangle$ Homeo reaction $v_i: X+X \rightarrow k_i P_i$ add $2;P_i$ and 1;0 to $\langle X, v_i \rangle$

2: Read the result by rows:

$$A = ?b_{(r)}; (A|A) \oplus ?c_{(r)}; (A|B) \oplus !c_{(r)}; 0$$
$$B = t_{(s)}; A \oplus !b_{(r)}; 0$$

	$x_{(s)}$	b _(r)	C _(r)
A		?;A A	?;A B !;0
В	τ; Α	!;0	



- x: B →^s A b: A+B →^r A+A c: A+A →^{2r} A+B
- 1: Fill the matrix by columns:

Degradation reaction $v_i: X \rightarrow k_i P_i$ add $t;P_i$ to $\langle X, v_{ii} \rangle$. Hetero reaction $v_i: X+Y \rightarrow k_i P_i$ add $2;P_i$ to $\langle X, v_i \rangle$ and 1;0 to $\langle Y, v_i \rangle$ Homeo reaction $v_i: X+X \rightarrow k_i P_i$ add $2;P_i$ and 1;0 to $\langle X, v_i \rangle$

2: Read the result by rows:

$$A = ?b_{(r)}; A \oplus ?c_{(r)}; (A|B) \oplus !c_{(r)}; 0$$
$$B = t_{(s)}; A \oplus !b_{(r)}; A$$





- x: B →^s A b: A+B →^r A+A c: A+A →^{2r} A+B
- 1: Fill the matrix by columns:

Degradation reaction $v_i: X \rightarrow k_i P_i$ add $t;P_i$ to $\langle X, v_{ii} \rangle$. Hetero reaction $v_i: X+Y \rightarrow k_i P_i$ add $2;P_i$ to $\langle X, v_i \rangle$ and 1;0 to $\langle Y, v_i \rangle$ Homeo reaction $v_i: X+X \rightarrow k_i P_i$ add $2;P_i$ and 1;0 to $\langle X, v_i \rangle$

2: Read the result by rows:

$$A = ?b_{(r)}; A \oplus ?c_{(r)}; B \oplus !c_{(r)}; A$$
$$B = t_{(s)}; A \oplus !b_{(r)}; A$$





From Chemistry to CGF: Pi(C)



From uniquely-labeled (v:) chemical reactions C to a CGF Pi(C):





Discrete-State Semantics



Discrete Semantics of Reactions



Discrete Semantics of Reagents



CTMC



Discrete State Equivalence

- Def: 🗯 is equivalent CTMC's (isomorphic graphs with same rates).
- Thm: E 🗯 Ch(E)
- Thm: C = Pi(C)



- For each E there is an E' \approx E that is detangled (E' = Pi(Ch(E)))

Interacting Automata = Discrete Chemistry

This is enough to establish that the process algebra is really faithful to the chemistry.

But CTMC are not the "ultimate semantics" because there are still questions of when two different CTMCs are actually equivalent (e.g. "lumping").

The "ultimate semantics" of chemistry is the *Chemical Master Equation* (derivable from the Chapman-Kolmogorov equation of the CTMC).



Entangled vs Detangled





Entangled: Two reactions on one channel Detangled: Two reactions on two separate channels

We need a semantics of automata that identifies automata that have the "same chemistry". No process algebra equivalence is like this!

Entangled automata lead to more compact models than in chemistry.

Detangled automata are in simple correspondence with chemistry.

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From Discrete to Continuous Chemistry

The "Type System" of Chemistry

The International System of Units (SI) defines the following physical units, with related derived units and constants; note that *amount of substance* is a base unit in SI, like length and time:

mol	(a base unit)	mole, unit of amount of substance
т	(a base unit)	meter, unit of <i>length</i>
5	(a base unit)	second, unit of time
L = 0.0)01 <i>.m</i> ³	liter (volume)
M = m	ol-L ⁻¹	molarity (concentration of substance)
N _A :m	$ol^{-1} \cong 6.022 \times 10^{23}$	Avogadro's number (number of particles per amount of substance)

For a substance X:mol, we write [X]:M for the concentration of X, and $[X]^{\bullet}:M \cdot s^{-1}$ for the time derivative of the concentration.

A continuous chemical system (C,V) is a system of chemical reactions C plus a vector of initial concentrations V_X : M, one for each species X.

The rates of unary reactions have dimension s^{-1} .

The rates of binary reactions have dimension $M^{-1}s^{-1}$.

(because in both cases the rhs of an ODE should have dimension $M \cdot s^{-1}$).

For a given volume of solution V, the volumetric factor γ of dimension M⁻¹ is:

 $\gamma: M^{-1} = N_A V$ where $N_A: mol^{-1}$ and V: L

 $\#X / \gamma : M =$ concentration of X molecules $\gamma [X] : 1 =$ total number of X molecules (rounded to an integer).

The Gillespie Conversion

Discrete Chemistry	Continuous Chemistry	$\gamma = N_A V$:M ⁻¹
initial quantities #A ₀	initial concentration [A] ₀	$ \text{with } [A]_0 = \#_A$	Α ₀ /γ
A→r A'	$A \to^k A'$	with <mark>k = r</mark>	:s ⁻¹
A+B ⊶• A'+B'	$A + B \rightarrow^k A' + B'$	with <mark>k = rγ</mark>	:M ⁻¹ s ⁻¹
A+A ⊶•r A'+A″	$A+A \rightarrow^k A'+A''$	with <mark>k = rγ/2</mark>	:M ⁻¹ s ⁻¹

V = interaction volume N_A = Avogadro's number

Think $\gamma = 1$ i.e. V = 1/N_A

M = mol·L⁻¹ molarity (concentration)



$Cont_{\gamma}$ and $Disc_{\gamma}$

4.2-3 Definition: Cont_y and Disc_y

For a volumetric factor $\gamma:M^{-1}$, we define a translation $Cont_{\gamma}$ from a discrete chemical systems (C,P), with species X and initial molecule count $\#X_0 = \#X(P)$, to a continuous chemical systems (C,V) with initial concentration $[X]_0 = V_X$. The translation $Disc_{\gamma}$ is its inverse, up to a rounding error $\lceil \gamma[X]_0 \rceil$ in converting concentrations to molecule counts. Since γ is a global conversion constant, we later usually omit it as a subscript.

$Cont_{\gamma}(X \rightarrow^{r} P)$	$= X \rightarrow^k P$	with $k = r$,	r:s ⁻¹	k:s ⁻¹
$Cont_{\gamma}(X+Y \rightarrow^{r} P)$	$= X+Y \rightarrow^{k} P$	with $\mathbf{k} = \mathbf{r} \boldsymbol{\gamma}$	r:s ⁻¹	k:M ⁻¹ s ⁻¹
$Cont_{\gamma}(X+X \rightarrow^{r} P)$	$= X + X \rightarrow^{k} P$	with $\mathbf{k} = \mathbf{r}\gamma/2$	r:s ⁻¹	k:M ⁻¹ s ⁻¹
$Cont_{\gamma}(\#X_0)$	= [X] ₀	with $[X]_0 = #X_0/\gamma$	$X_0:mol$	$[X]_0:M$
$Disc_{\gamma}(X \rightarrow^{k} P)$	$= X \rightarrow^{r} P$	with $r = k$,	k:s ⁻¹	r:s ⁻¹
$Disc_{\gamma}(X \to^{k} P)$ $Disc_{\gamma}(X+Y \to^{k} P)$	$= X \rightarrow^{r} P$ $= X+Y \rightarrow^{r} P$	with $r = k$, with $r = k/\gamma$	k:s ⁻¹ k:M ⁻¹ s ⁻¹	r:s ⁻¹ r:s ⁻¹
$Disc_{\gamma}(X \rightarrow^{k} P)$ $Disc_{\gamma}(X+Y \rightarrow^{k} P)$ $Disc_{\gamma}(X+X \rightarrow^{k} P)$	$= X \rightarrow^{r} P$ $= X+Y \rightarrow^{r} P$ $= X+X \rightarrow^{r} P$	with $r = k$, with $r = k/\gamma$ with $r = 2k/\gamma$	k:s ⁻¹ k:M ⁻¹ s ⁻¹ k:M ⁻¹ s ⁻¹	r:s ⁻¹ r:s ⁻¹ r:s ⁻¹

 $Ch_{\gamma} := Cont_{\gamma} \circ Ch$



Continuous-State Semantics (summary)



Same Semantics

Could chemistry itself be that semantics? No: different sets of reactions can have the same behavior!



From Reactions to ODEs (Law of Mass Action)



Continuous State Equivalence

• Def: \approx is equivalence of polynomials over the field of reals.



- For each E there is an $E' \approx E$ that is detangled (E' = Pi(Ch(E)))
- For each E in automata form there is an an E' ≈ E that is detangled and in automata form (E' = Detangle(E)).

GMA ≠ CME



$A+A \rightarrow^{2r} A =? A+A \rightarrow^{r} 0$



(For conservation of mass, consider instead $A+A \rightarrow^{2r} A+B$ vs. $A+A \rightarrow^{r} B+B$)

Continuous vs. Discrete Groupies



directive sample 5.0 1000 directive plot B0; A0	directive sample 5.0 1000 directive plot B(); A()	directive sample 5.0 1000 directive plot B(j: A()		
new a@1.0:chan() new b@1.0:chan()	new a@1.0:chan() new b@1.0:chan()	new a@1.0:chan() new b@1.0:chan()		Groupe ODEs - Groupies Hysteric 1,mat
et A0 = do !a; A0 or ?b; B0 nd B0 = do !b; B0 or ?a; A0	let A() = do la; A() or 2b; 2b; B() and B() = do lb; B() or 2a; 2a; A()	let A() = do !a; A() or ?b; ?b; ?b; B() and B() = do !b; B() or ?a; ?a; ?a; A()	(0:0:0015:0) r=10 kt-20 d dv1(d+dv1-2) 2000 0	[0:0.001:5.0] r=1.0 k=1.0 A dx1/dt=x1*x4+x3*x1-x1+x4, 2000.0 A' dx2/dt=x3*x1-x3*x2+x1-x2, 0.0
ld() = la; Ad() Bd() = lb; Bd()	let Ad() = la: Ad() and Bd() = lb; Bd()	let Ad() =!a; Ad() and Bd() = fb; Bd()	B dx2/dt = (x1-x2), 0.0	B dx3/dt=x3*x2-x1*x3-x3+x2, 0.0 B' dx4/dt=x1*x3-x1*x4+x3-x4, 0.0
2000 of A0 1 of (Ad() Bd())	run 2000 of A() run 1 of (Ad() Bd())	run 2000 of A() run 1 of (Ad() Bd())		

Scientific Predictions



After a while, all 4 states are almost equally occupied.

The 4 states are almost never equally occupied.

Discrete Analysis Techniques

The Program vs. the State Space

The "program":





Potentially infinite

Simulation

- Run "the program" through a walk in states space.
- Basic stochastic algorithm: Gillespie
 - Exact (i.e. based on physics) stochastic simulation of chemical kinetics.
 - Can compute concentrations and reaction times for biochemical networks.
- Stochastic Process Algebras
 - Now many [BioSPi, SPiM, BioPEPA, BetaBinders, ...]
- Hybrid approaches
 - Continuous + discrete/stochastic switching


Control Flow Analysis

- Who may call who?
 - Overapproximation of behavior used to answer questions about what "cannot happen".

What event may (or may not) have been involved in reaching this state?



Causality Analysis

- What event caused what other event or state to happen?
 - E.g.: if in all possible executions one event always precedes another.

 Need a different level of representation (the "event space")

- Petri Nets
- Event Structures





Abstract Interpretation

• Precisely relating abstract views to more concrete views of the system



Modelchecking

- Asking questions (in Temporal Logic) about structure of a (finite) state space.
- Various flavors of modelchecking:
 - Temporal
 - About paths through state space
 - Quantitative
 - About quantitative measures of states
 - Probabilistyc/Stochastic
 - About probabilities of reaching states.

Is this state a necessary checkpoint to reach this state?

Bisimulation

- Are two programs generating the same state space?
 - E.g.: Is a compact description of a system equivalent to a more detailed one in all possible environments?



Conclusions

Conclusions

- Process Algebra
 - An extension of automata theory to populations of interacting automata
 - Modeling the behavior of individuals in an arbitrary environment
 - Compositionality (combining models by juxtaposition)
- Connections between modeling approaches
 - Connecting the discrete/concurrent/stochastic/molecular approach
 - to the continuous/sequential/deterministic/population approach
- Connecting syntax with semantics
 - Syntax = model presentation (equations/programs/diagrams/blobs etc.)
 - Semantics = state space (generated by the syntax)
- Ultimately, connections between analysis techniques
 - We need (and sometimes have) good semantic techniques to analyze state spaces (e.g. calculus, but also increasingly modelchecking)
 - But we need equally good syntactic techniques to structure complex models (e.g. compositionality) and analyze them (e.g. process algebra)
- A bright future for Computer Science and Logic in modern Biology
 - Biology needs good analysis techniques for discrete systems analysis (modal logics, modelchecking, causality analysis, abstract interpretation, ...)

